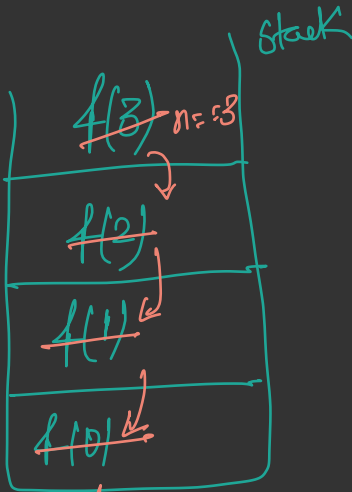
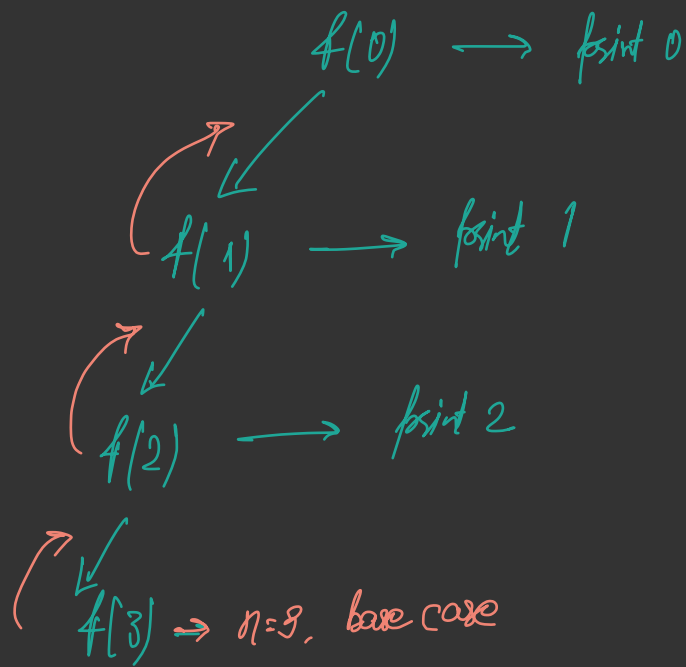


Recursion  $\rightarrow$  when a function calls itself until a specified cond is met

```

1 function f(n) {
2   if(n == 3) {
3     return;
4   }
5   console.log(n) //?  $\rightarrow$  0, 1, 2
6   n++;
7   f(n);
8 }
9
10 function main() {
11   f(0);
12 }
13
14 main(); //?
  
```

$f(n)$



control goes back to main  $f^n$

lec2 Problems on recursion

- $\rightarrow$  print name 5 times
- $\rightarrow$  print linearly from 1 to N
- $\rightarrow$  print from N to 1
- $\rightarrow$  print linearly from 1 to N (but by backtracking)
- $\rightarrow$  print from N to 1 (by back-track)

print name 5 times.

```

function main() {
  n = 5;
  f(1, n);
}

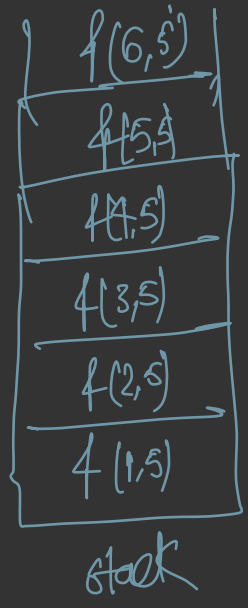
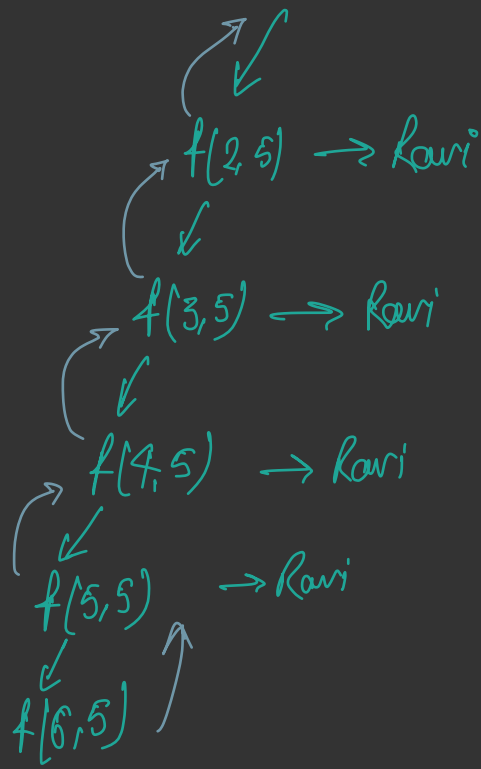
```

```

function f(i, n) {
  if (i > n)
    return;
  print("Ravi");
  f(i + 1, n);
}

```

$f(1, 5) \rightarrow \text{Ravi}$



Print 1 to N

```

main() {
  n = 5;
  f(1, n);
}

```

```

f(i, n) {
  if (i > n) return;
  print(i);
  f(i + 1, n);
}

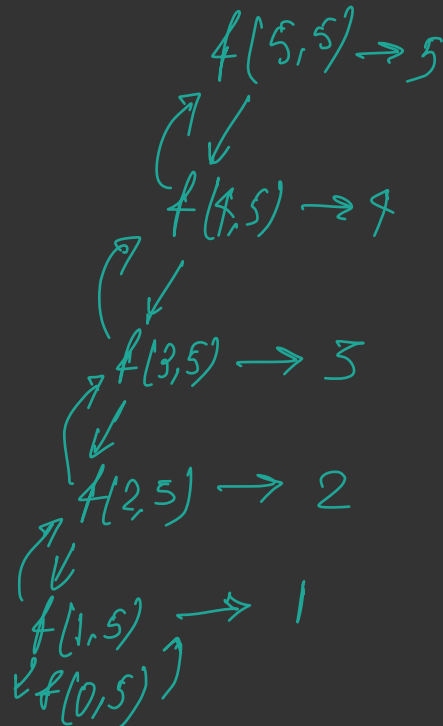
```

print N to 1

```

main() {
  n = 5;
  f(n, n);
}
f(i, n) {
  if (i < 1) return;
  print(i);
  f(i - 1, n);
}

```



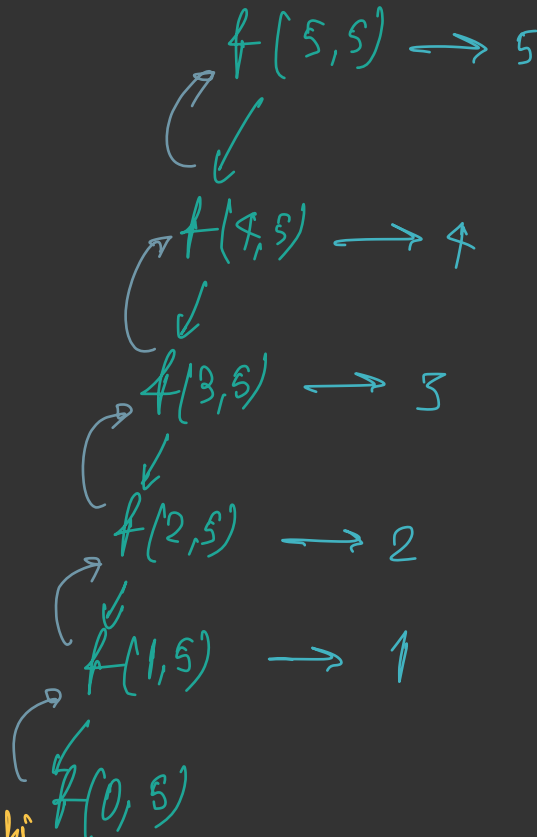
print linearly from 1 to N backtracking

do not use  $i+1$ ,  
but can use  $i-1$

recursion first  
then action

```
main() {
    n = 5
    f(n, n)
}
```

```
f(i, n) {
    if (i < 1) return;
    f(i-1, n);
    print(i)
}
```

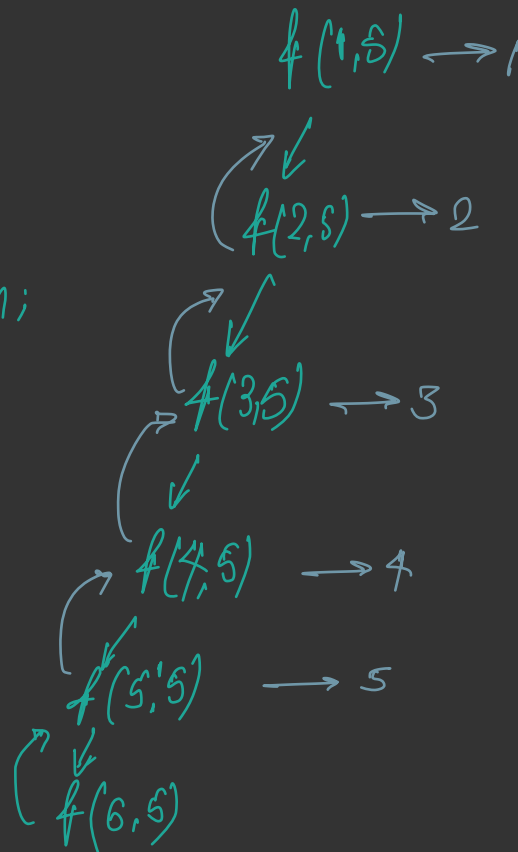


1, 2, 3, 4, 5

print N to 1 using backtracking

```
main() {
    n = 5
    f(1, n)
}
```

```
f(i, n) {
    if (i > n) return;
    f(i+1, n);
    print(i)
}
```



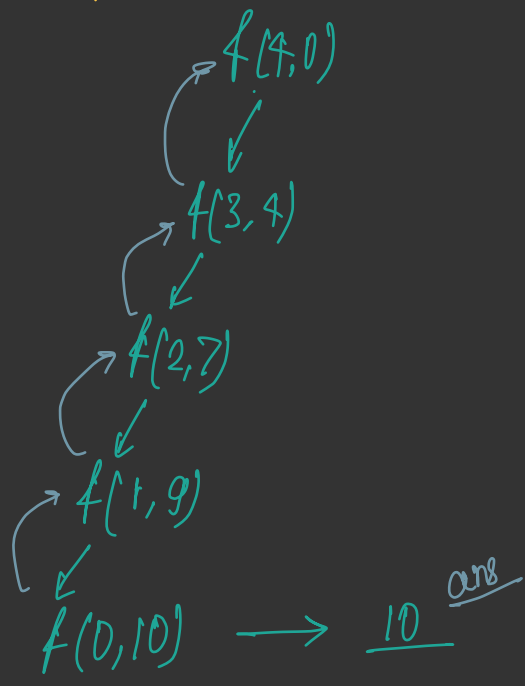
5 4 3 2 1

# lec 3 Parameterised and functional recursion

Q. sum of first N numbers.

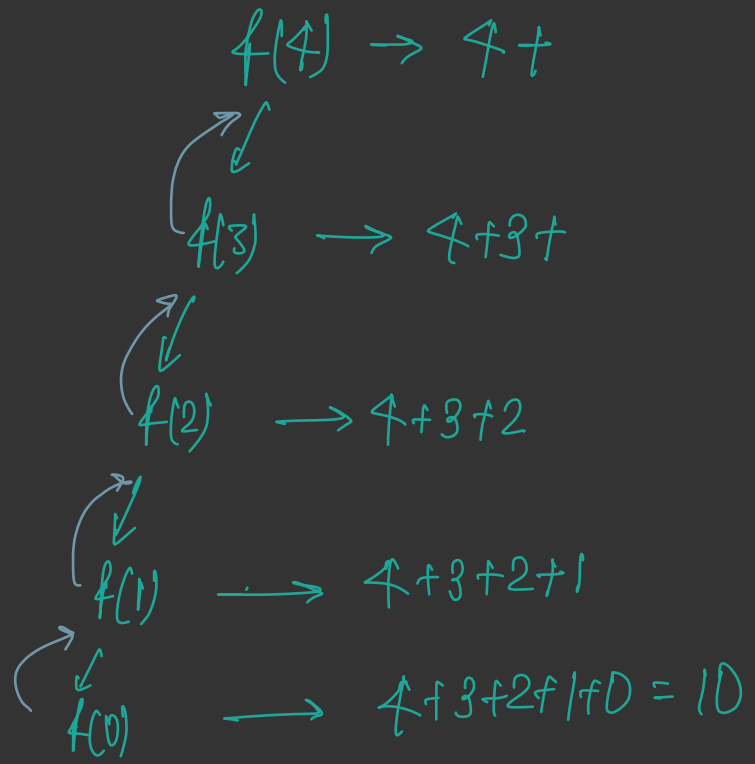
parameter ↙ functional ↘

```
main() {
    n = 4
    f(n, 0)
}
f(i, sum) {
    if (i < 1) {
        print(sum)
        return;
    }
    f(i-1, sum+i)
}
```



functional.

```
main() {
    n = 4
    f(n)
}
f(i) {
    if (i == 0) return 0;
    return i + f(i-1)
}
```



TC → O(N)

SC → O(N)

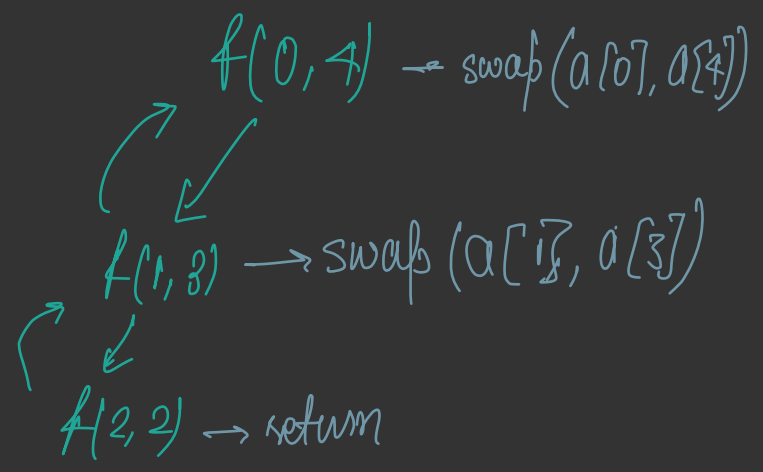
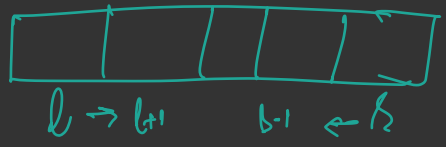
↗ stack



# lec 4. problems on functional recursion

## Q. reverse an array.

$a = 1, 2, 3, 4, 5$       $\text{swap}(i, j)$       $n = 5$   
 $a = 5, 4, 3, 2, 1$



```

f(l, s) {
    if (l >= s) return;
    swap(a[l], a[s]);
    f(l+1, s-1);
}
    
```

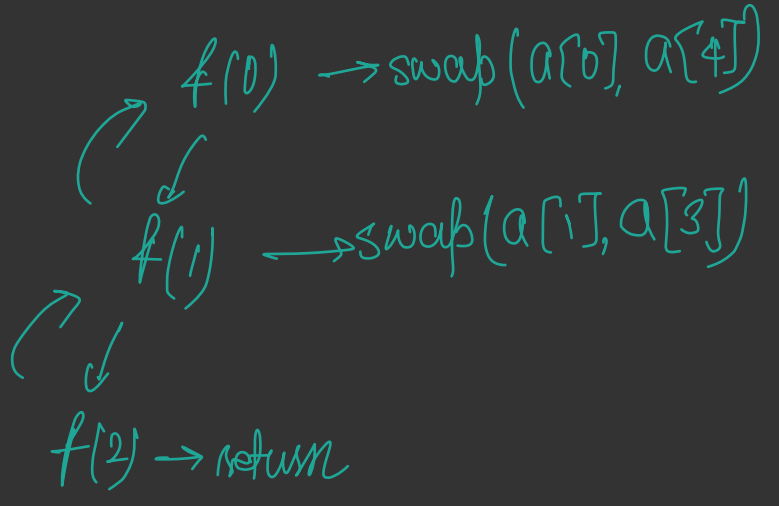
```

main() {
    l = 0;
    s = n-1;
    f(l, s);
}
    
```

using one parameter  
 $l$ ,  $s = n-1-l$   
 if ( $l \geq n/2$ ) return;

```

f(i) {
    if (i >= n/2) return;
    swap(a[i], a[n-i-1]);
    f(i+1);
}
    
```



$$n/2 = 5/2 = 2$$

```

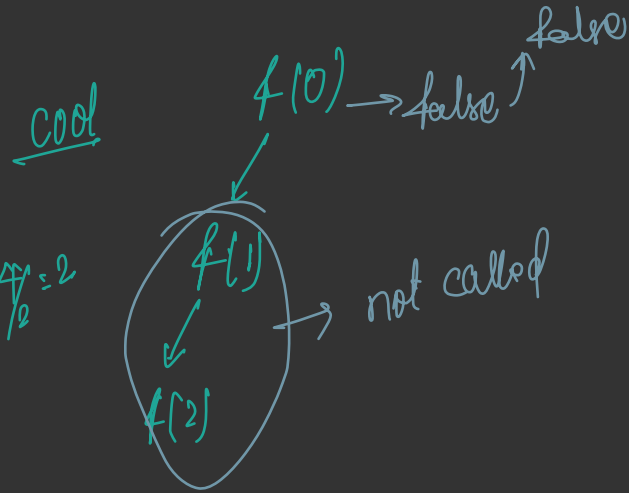
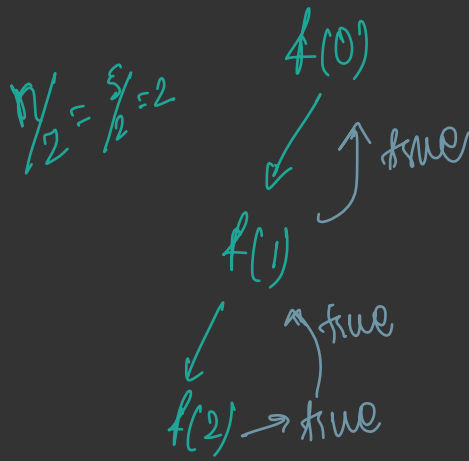
main() {
    n = 5;
    f(0);
}
    
```

Q check if given string is palindhrome

```

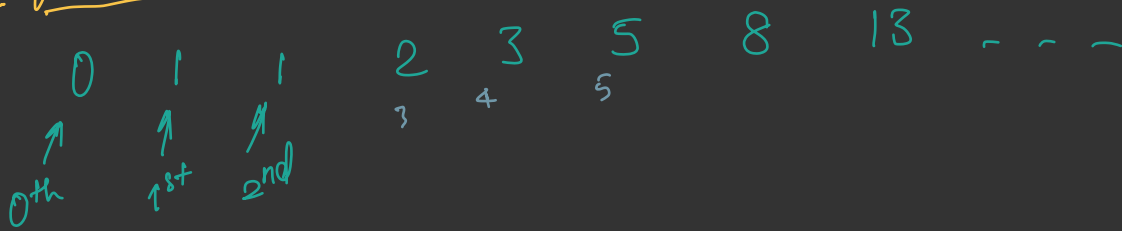
main() {
    s = "cavac"
    f(0)
}
f(i) {
    if (i >= n/2) return true;
    if (s[i] != s[n-1-i])
        return false;
    return f(i+1)
}
    
```

TC  $\rightarrow O(N/2)$   
 SC  $\rightarrow O(N/2)$



lec 5  $\rightarrow$  multiple recursion calls.

Q. fibonacci numbers



$N \rightarrow f(n) \rightarrow$  nth fibonacci num

$f(3) \rightarrow 3^{\text{rd}} \rightarrow 2$

$f(2) \rightarrow 2^{\text{nd}} \rightarrow 1$

$f(4) \rightarrow 4^{\text{th}} \rightarrow 3$

$f(5) \rightarrow f(4) + f(3) = 3 + 2 = 5$

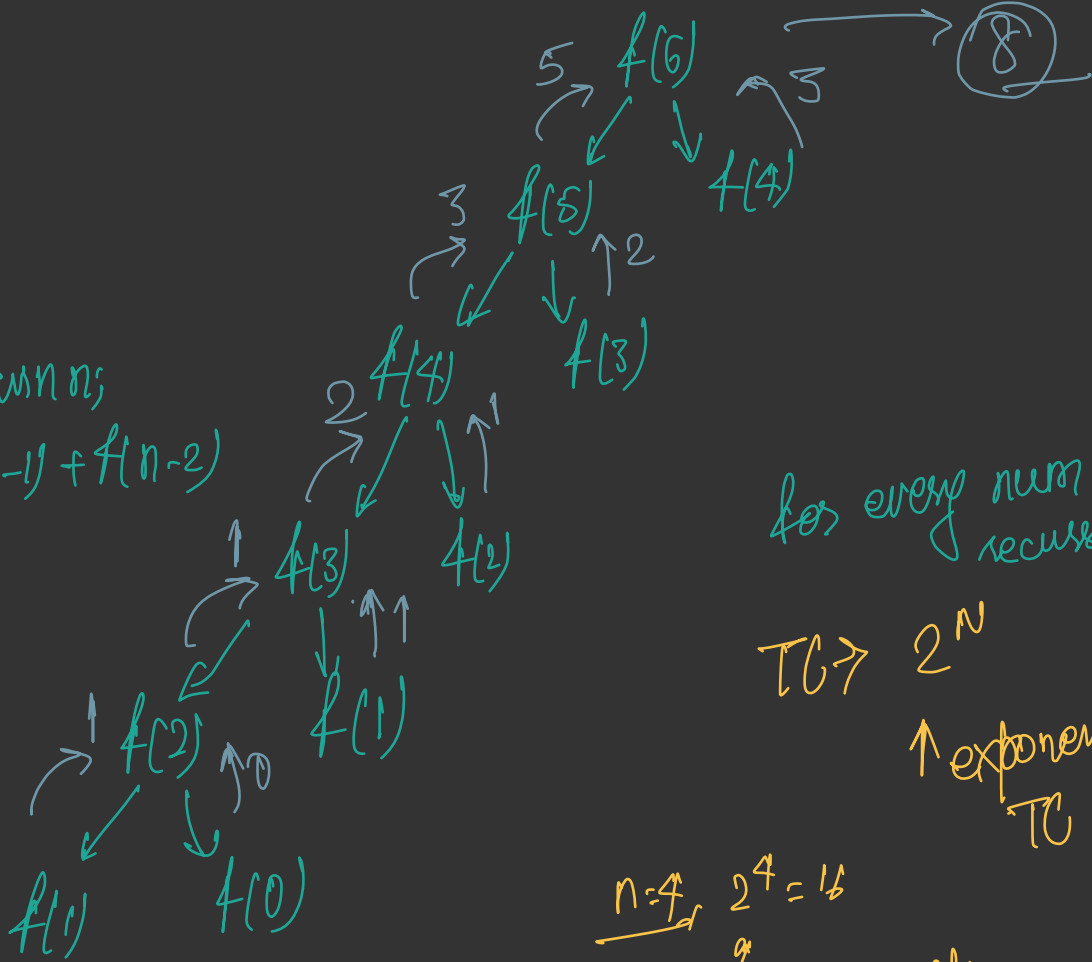
$f(0) = 0$

$f(1) = 0$

$f(2) = f(0) + f(1) = 0 + 1 = 1$

```
main() {
  f(6)
}
```

```
f(n) {
  if (n <= 1) return n;
  return f(n-1) + f(n-2)
}
```



for every num 2 recursion call

TC  $\rightarrow 2^N$

$\uparrow$  exponential TC

$n=4, 2^4 = 16$

actual 9 calls  
so it's exactly  $2^N$  but it's exponential

lec 6. Recursion on subsequences

$\rightarrow$  a contiguous / non-contiguous sequence which follows the order of array

{ 3, 1, 2 }

$\rightarrow$  { 3, 2 }

[ 1, 2 ]

[ ]

Q. print all subsequences.

take / not Take

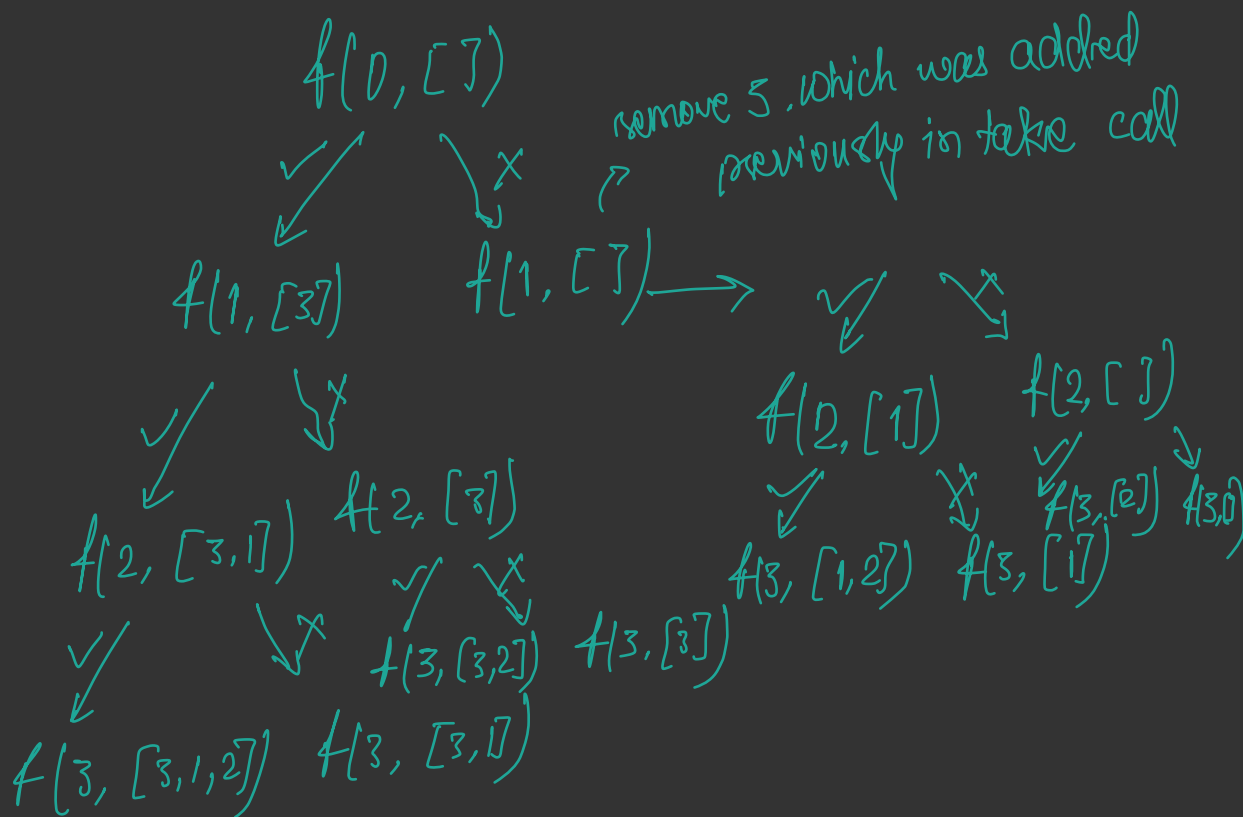
$f(ind, [ ])$

```

if (ind >= n)
    return [ ]
[ ].add(nums[ind]);
f(ind+1, [ ]) → take
[ ].remove(nums[ind]);
f(ind+1, [ ]) → not take
}

```

$[3, 1, 2]$



$[3, 1, 2], [3, 1], [3, 2], [3], [1, 2], [1], [2], [ ]$

TC  $\rightarrow O(2^N \times N)$

SG  $\rightarrow O(N)$

```

1 function f(i, n, nums, dp) {
2   if (i >= n) {
3     return dp;
4   }
5
6   let withCurr = f(i + 1, n, nums, [...dp, nums[i]]);
7   let withoutCurr = f(i + 1, n, nums, dp);
8
9   return [...withCurr, ...withoutCurr];
10 }
11
12 function printSubSequences(nums) {
13   let n = nums.length;
14   let dp = [];
15   return f(0, n, nums, dp);
16 }

```

```

function f(i, n, nums, dp){
  if(i >= n){
    console.log(dp); [ 3, 1, 2 ],
    return [ 3, 1 ],
  }
  dp.push([ 3, 2 ])
  dp.push([ 3 ])
  f(i+1, n, nums, [...dp, nums[i]]);
  dp.pop();
  f(i+1, n, nums, dp);
}

```

```

function printSubSequences(nums){
  let n = nums.length;
  let dp = [];
  f(0, n, nums, dp);
}

const nums = [3,1,2];
printSubSequences(nums); undefined

```

```

1 function f(i, n, nums, dp){
2   if(i >= n){
3     console.log(dp);
4     return;
5   }
6
7   dp.push(nums[i]);
8   f(i+1, n, nums, dp);
9   dp.pop();
10  f(i+1, n, nums, dp);
11 }
12
13 function printSubSequences(nums){
14   let n = nums.length;
15   let dp = [];
16   f(0, n, nums, dp);
17 }
18
19 const nums = [3,1,2];
20 printSubSequences(nums);

```

lec 7 → all kind of patterns in recursion

→ take & not-take

→ pick & not-pick

Q. find subarray which sum is k

[1, 2, 1], k=3

f(i, ds, s) {

if (i == n) {

if (s == sum)

print(ds)

return

}

ds.add(arr[i])

s += arr[i]

f(i+1, ds, s) → pick

ds.remove(arr[i])

s -= arr[i]

f(i+1, ds, s) → not pick

}

```

void printS(int ind, vector<int> &ds, int s, int sum, int arr[], int n) {
    if(ind == n) {
        if(s == sum) {
            for(auto it : ds) cout << it << " ";
            cout << endl;
        }
        return;
    }

    ds.push_back(arr[ind]);
    s += arr[ind];

    printS(ind+1, ds, s, sum, arr, n);

    s -= arr[ind];
    ds.pop_back();

    // not pick
    printS(ind+1, ds, s, sum, arr, n);
}

int main() {
    #ifndef ONLINE_JUDGE
    freopen("input.txt", "r", stdin);
    freopen("output.txt", "w", stdout);
    #endif
    int arr[] = {1, 2, 1};
    int n = 3;
    int sum = 2;
    vector<int> ds;
    printS(0, ds, 0, sum, arr, n);
}

```

Q. find any one sub-sequence whose sum is k

⇒ technique to print only one answer  
in base case if condition satisfied return true

```

bool flag = false;
void printS(int ind, vector<int> &ds, int s, int sum, int arr[], int n) {
    if(ind == n) {
        if(s == sum && flag == false) {
            flag = true;
            for(auto it : ds) cout << it << " ";
            cout << endl;
        }
        return;
    }

    ds.push_back(arr[ind]);
    s += arr[ind];

    printS(ind+1, ds, s, sum, arr, n);

    s -= arr[ind];
    ds.pop_back();

    // not pick
    printS(ind+1, ds, s, sum, arr, n);
}

```

not preferred

base case →  
cond → satisfied  
return true

else  
return false

if (f(i) == true)  
return true

f(n)  
→ return false

```

bool printS(int ind, vector<int> &ds, int s, int sum, int arr[], int n) {
    if(ind == n) {
        // condition satisfied
        if(s == sum) {
            for(auto it : ds) cout << it << " "; → print
            cout << endl;
            return true;
        }
        // condition not satisfied
        else return false;
    }

    ds.push_back(arr[ind]);
    s += arr[ind];

    if(printS(ind+1, ds, s, sum, arr, n) == true) {
        return true;
    }

    s -= arr[ind];
    ds.pop_back();

    // not pick
    if(printS(ind+1, ds, s, sum, arr, n) == true) return true;

    return false;
}

```

✓ TC →

Q. count subsequences with sum k

[1, 2, 1] k=2

count

base case

return 1 → cond<sup>n</sup> satisfied

return 0 → cond<sup>n</sup> not satisfied

left = f()

right = f()

return (left + right)

TC →  $O(2^N)$

SC → 0

```

1 function f(i, n, nums, k) {
2     if (i >= n) {
3         return k == 0 ? 1 : 0;
4     }
5     if(k == 0) return 1;
6
7     let left = f(i+1, n, nums, k-nums[i]);
8     let right = f(i+1, n, nums, k)
9     return left + right;
10 }
11
12 function countSubsequences(nums, k) {
13     let n = nums.length;
14     return f(0, n, nums, k);
15 }
16
17 const nums = [1, 2, 1];
18 console.log(countSubsequences(nums, 2));

```

lec8. merge sort → already completed in merge sort notes

lec9. quick sort → already completed in sorting notes

lec10. combination sum

pick, not pick  
& infinite supplies

$a = [2, 3, 6, 7]$ ,  $k = 7$

### 39. Combination Sum

Medium

15.7K

312

☆

🔄

Airbnb

Adobe

Microsoft

Given an array of **distinct integers** `candidates` and a target integer `target`, return a **list of all unique combinations** of `candidates` where the chosen numbers sum to `target`. You may return the combinations in **any order**.

The **same number** may be chosen from `candidates` an **unlimited number of times**. Two combinations are unique if the **frequency** of at least one of the chosen numbers is different.

The test cases are generated such that the number of unique combinations that sum up to `target` is less than **150** combinations for the given input.

#### Example 1:

**Input:** candidates = [2,3,6,7], target = 7

**Output:** [[2,2,3],[7]]

#### Explanation:

2 and 3 are candidates, and  $2 + 2 + 3 = 7$ . Note that 2 can be used multiple times.

7 is a candidate, and  $7 = 7$ .

These are the only two combinations.

$f(\text{ind}, \text{target}, \text{ds})$

✓  
ds.add(a[ind])

✗  
ds.pop(a[i])  
 $f(\text{ind}+1, \text{target}, \text{ds})$

←  $f(\text{ind}, \text{target} - a[\text{ind}], \text{ds})$

$\text{if}(a[\text{ind}] \leq \text{target})$

TC →  $2^t \times k$

exponential

SC →  $k \times t$

base case

$\text{if}(\text{ind} == n)$   
 $\text{if}(\text{target} == 0)$  log(ds)  
else return

```
1 var combinationSum = function (candidates, target) {
2   let ds = [];
3   let n = candidates.length;
4   let ans = [];
5   f(0, n, target, candidates, ds, ans);
6   return ans;
7 };
8
9 function f(i, n, k, a, ds, ans) {
10  if (i == n) {
11    if (k == 0) ans.push([...ds]);
12    return;
13  }
14
15  if (a[i] <= k) {
16    ds.push(a[i]);
17    f(i, n, k - a[i], a, ds, ans);
18    ds.pop();
19  }
20
21  f(i + 1, n, k, a, ds, ans);
22 }
```



# lec 11 combination sum 2

```

1 var combinationSum2 = function (candidates, target) {
2   let n = candidates.length;
3   let ds = [];
4   let ans = new Set();
5   f(0, n, target, candidates, ds, ans);
6   return Array.from(ans, (seq) => seq.split(",").map(Number));
7 };
8
9 function f(i, n, k, a, ds, ans) {
10  if (i == n) {
11    if (k == 0) {
12      ans.add(
13        ds
14          .slice()
15          .sort((a, b) => a - b)
16          .toString()
17      );
18    }
19    return;
20  }
21
22  if (a[i] <= k) {
23    ds.push(a[i]);
24    f(i + 1, n, k - a[i], a, ds, ans);
25    ds.pop();
26  }
27  f(i + 1, n, k, a, ds, ans);
28 }

```

TLE

## 40. Combination Sum II

Medium 8.3K 209 ☆

Given a collection of candidate numbers (`candidates`) and a target number (`target`), find all unique combinations in `candidates` where the candidate numbers sum to `target`.

Each number in `candidates` may only be used **once** in the combination.

**Note:** The solution set must not contain duplicate combinations.

### Example 1:

```

Input: candidates = [10,1,2,7,6,1,5], target = 8
Output:
[
  [1,1,6],
  [1,2,5],
  [1,7],
  [2,6]
]

```

TC  $\rightarrow O(2^t \times k \log n)$   
 $\downarrow$  set

sort the given array

avg length of subsequence

TC  $\rightarrow 2^N \times k$

f(ind, target, ds, ans) {

if (target == 0)  
 ans.push(ds)  
 return

for (let i = ind; i < n; i++) {  
 if (i > ind && a[i] == a[i-1]) continue;  
 if (a[i] > target) break;

ds.push(a[i])  
 f(i+1, a, target - a[i], ans, ds);  
 ds.pop()

}

```

1 var combinationSum2 = function (candidates, target) {
2   let n = candidates.length;
3   let ds = [];
4   let ans = [];
5   candidates.sort((a, b) => a - b);
6   f(0, n, target, candidates, ds, ans);
7   return ans;
8 };
9
10 function f(ind, n, k, a, ds, ans) {
11   if (k == 0) {
12     ans.push([...ds]);
13     return;
14   }
15
16   for (let i = ind; i < n; i++) {
17     if (i > ind && a[i] == a[i - 1]) continue;
18     if (a[i] > k) break;
19
20     ds.push(a[i]);
21     f(i + 1, n, k - a[i], a, ds, ans);
22     ds.pop();
23   }
24 }

```



# lec 12. subset sum. $n=3$

num of subsets =  $2^n \Rightarrow 2^3 = 8$

- 3, 1, 2
- [ ]
- [ 3 ]
- [ 1 ]
- [ 2 ]
- [ 3, 1 ]
- [ 3, 2 ]
- [ 1, 2 ]
- [ 3, 1, 2 ]

M1 brute force → powerset  
 generate all subset  
 & traverse all subsets & return sum  
 TC →  $2^N \times N$

## M2 optimal

```

1 var subsetSum = function (nums) {
2   let n = nums.length;
3   let ans = [];
4   f(0, 0, n, nums, ans);
5   return ans.sort((a,b) => a-b);
6 };
7
8 function f(ind, sum, n, a, ans) {
9   if (ind == n) {
10    ans.push(sum);
11    return;
12  }
13
14  f(ind+1, sum, n, a, ans);
15  f(ind+1, sum + a[ind], n, a, ans)
16 }
17
18 const nums = [5,2,1]
19 subsetSum(nums); // 0,1,2,3,5,6,7,8

```

TC →  $2^N + 2^N \log(2^N)$  → size of the output  
 SC →

**Subset Sums** Basic Accuracy: 46.81% Submissions: 591 Points: 1

Given a list(Arr) of N integers, print sums of all subsets in it. **Output should be printed in increasing order of sums.**

**Example 1:**

**Input:**  
 N = 2  
 Arr = [2, 3]  
**Output:**  
 0 2 3 5  
**Explanation:**  
 When no elements is taken then Sum = 0.  
 When only 2 is taken then Sum = 2.  
 When only 3 is taken then Sum = 3.  
 When element 2 and 3 are taken then Sum = 2+3 = 5.

$f(0, 0)$       [3, 1, 2]  
 ↑ take, not take      ans  
 $f(i, sum)$  {  
 if (i == n) {  
   ans.push(sum)  
   return;  
 }  
 $f(ind+1, sum)$   
 $f(ind+1, sum + a[i])$   
 }

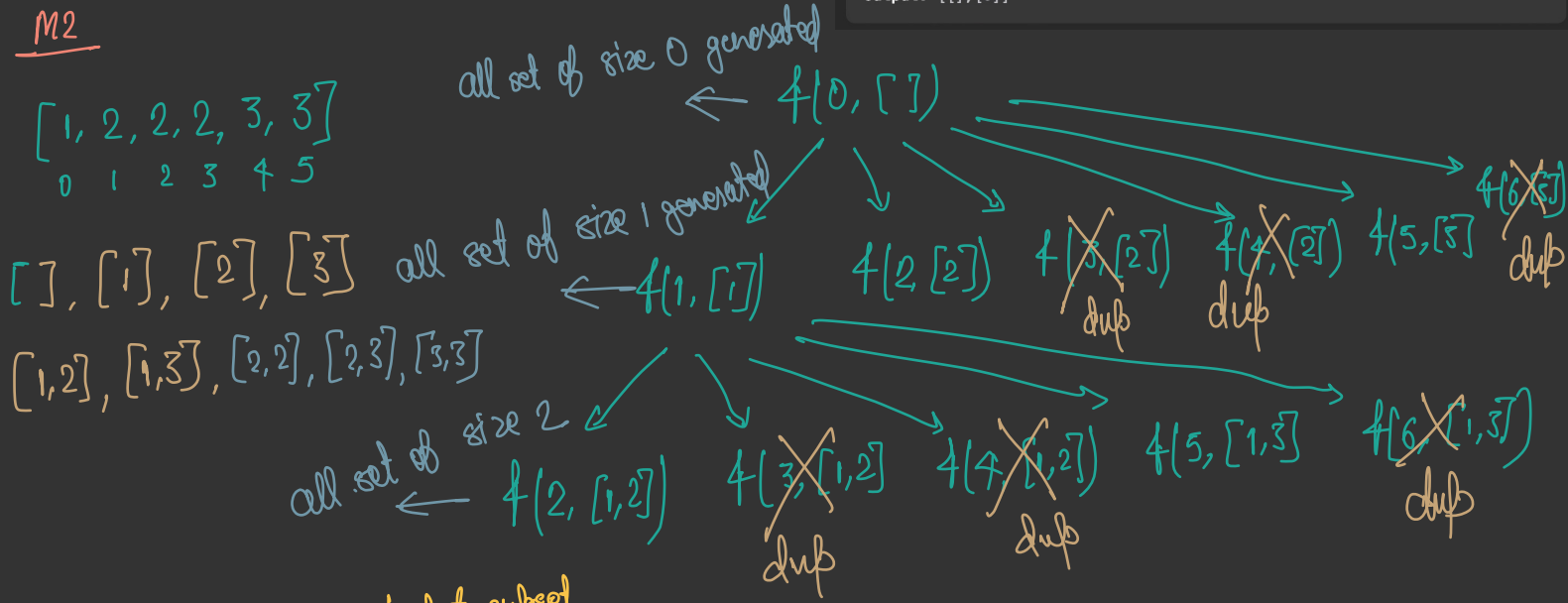
lec13 → subset sum 2

Mr. brute force

generate all subsets →  $2^N$   
 put these ( $2^N$ ) into a set  
 convert set to array of array

M2

[1, 2, 2, 2, 3, 3]  
 0 1 2 3 4 5



TC →  $2^N \times N$  ← to put subset into ans array

SC →  $O(2^N) \times O(K) + O(N)$   
 ↑ recursion

90. Subsets II

Medium 7.9K 226 ☆

Given an integer array `nums` that may contain duplicates, return all possible subsets (the power set).  
 The solution set must not contain duplicate subsets. Return the solution in any order.

Example 1:

Input: `nums = [1,2,2]`  
 Output: `[[], [1], [1,2], [1,2,2], [2], [2,2]]`

Example 2:

Input: `nums = [0]`  
 Output: `[[], [0]]`

```

f(0)
f(ind) {
    ans.push(ds)
    for (i = ind to n-1) {
        if (i > ind && nums[i] == nums[i-1]) continue;
        ds.push(nums[i]);
        f(i+1);
        ds.pop();
    }
}
    
```

```

1 function findSubsets(ind, nums, ds, ansList) {
2   ansList.push(...ds);
3   for (let i = ind; i < nums.length; i++) {
4     if (i !== ind && nums[i] === nums[i - 1]) continue;
5     ds.push(nums[i]);
6     findSubsets(i + 1, nums, ds, ansList);
7     ds.pop();
8   }
9 }
10
11 function subsetsWithDup(nums) {
12   nums.sort((a, b) => a - b);
13   let ansList = [];
14   findSubsets(0, nums, [], ansList);
15   return ansList;
16 }
17
18 const nums = [3, 1, 2];
19 subsetsWithDup(nums); //?
    
```

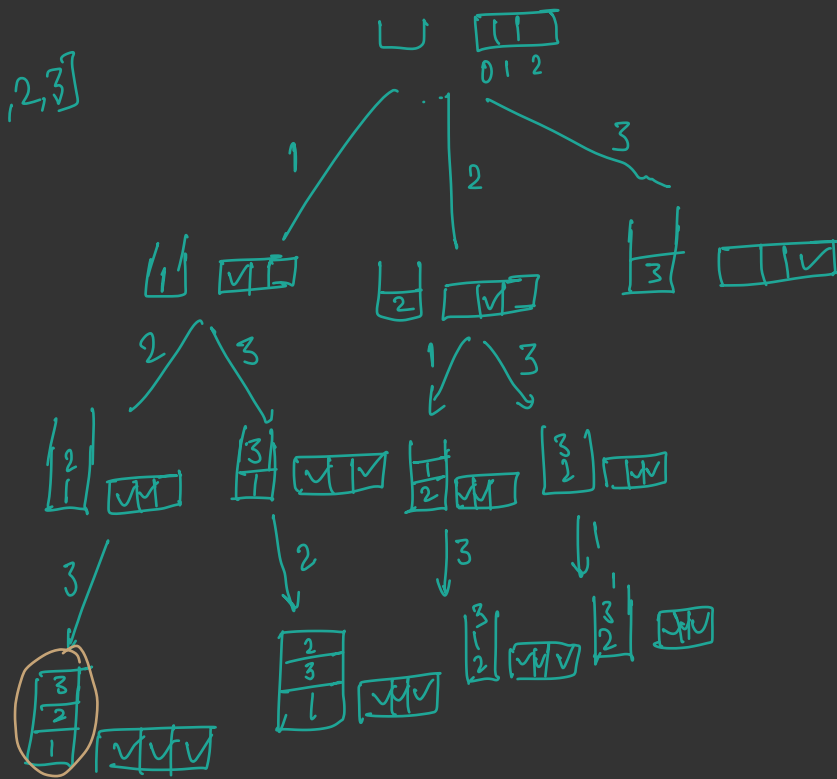
# leetcode print all permutations of a string/array

nums = [1, 2, 3] n=3

total permutation = 3!  
= 3 \* 2 \* 1  
= 6

- 1 2 3
- 1 3 2
- 2 1 3
- 2 3 1
- 3 1 2
- 3 2 1

[1, 2, 3]

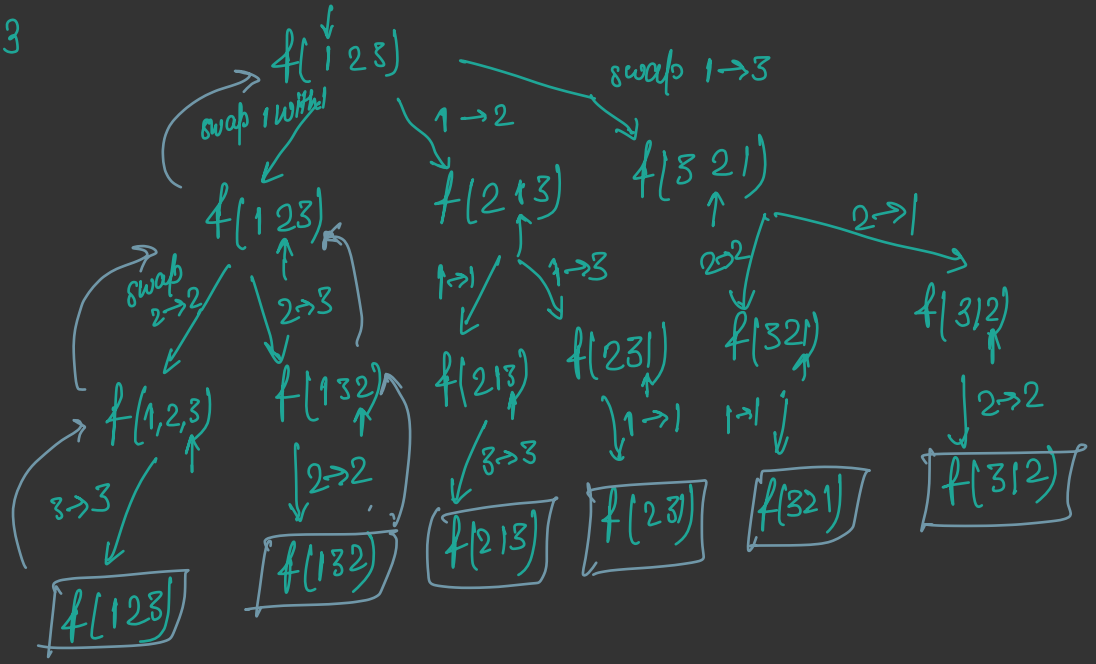


```
1 function f(nums, map, ans, ds){
2   if(ds.length == nums.length) {
3     ans.push([...ds]);
4     return;
5   }
6   for(let i=0; i<nums.length; i++){
7     if(!map[i]){
8       map[i] = true;
9       ds.push(nums[i]);
10      f(nums, map, ans, ds);
11      ds.pop();
12      map[i] = false;
13    }
14  }
15 }
16
17 function permutations(nums){
18   let ans = [];
19   let ds = [];
20   let n = nums.length;
21   let map = Array(n).fill(false);
22   f(nums, map, ans, ds);
23   return ans;
24 }
25
26 const nums = [3, 1, 2];
27 permutations(nums);
```

# lec 15 print all permutations of a string

[1, 2, 3] n=3

- 123
- 132
- 213
- 231
- 321
- 312



TC →  $N! \times N$

SC →  $O(N) + O(N!)$   
 ↑ recursion      ↑ ans

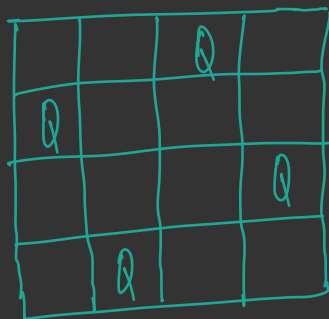
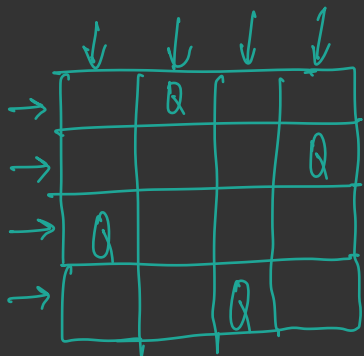
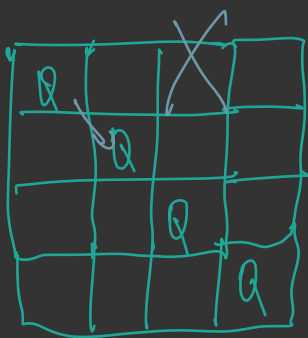
```

1 function f(ind, nums, ans){
2     if(ind == nums.length) {
3         let ds = [];
4         for(let i=0; i<nums.length; i++){
5             ds.push(nums[i]);
6         }
7         ans.push([...ds]);
8         return;
9     }
10
11     for(let i=ind; i<nums.length; i++){
12         [nums[i], nums[ind]] = [nums[ind], nums[i]];
13         f(ind+1, nums, ans);
14         [nums[ind], nums[i]] = [nums[i], nums[ind]];
15     }
16 }
17
18 function permutations(nums){
19     let ans = [];
20     f(0, nums, ans);
21     return ans;
22 }
23
24 const nums = [3, 1, 2];
25 permutations(nums); //?
    
```

# lect 6. N-Queens

- every row has 1 queen
- every col has 1 queen
- none of the Queen should attack each other

attack



## 51. N-Queens

Hard 9.9K 219

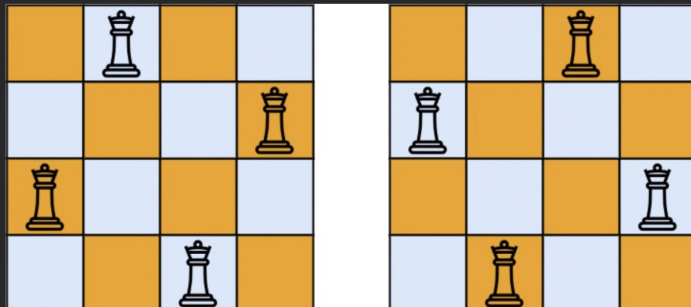
Amazon Adobe Microsoft

The **n-queens** puzzle is the problem of placing **n** queens on an **n x n** chessboard such that no two queens attack each other.

Given an integer **n**, return *all distinct solutions to the n-queens puzzle*. You may return the answer in **any order**.

Each solution contains a distinct board configuration of the n-queens' placement, where **'Q'** and **'.'** both indicate a queen and an empty space, respectively.

Example 1:

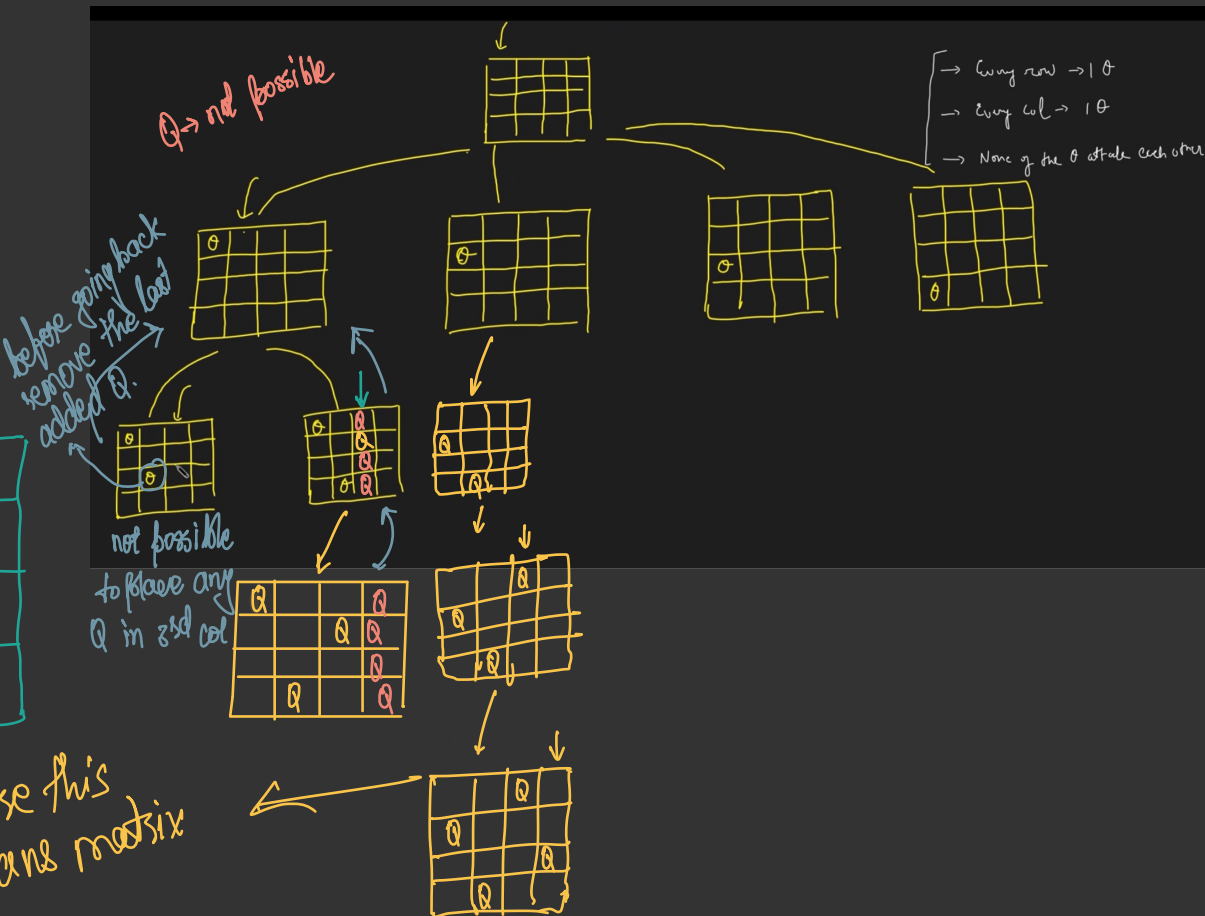


Input: n = 4

Output: [".Q...", "...Q", "Q...", "...Q."], [".Q.", "Q...", "...Q", ".Q.."]

Explanation: There exist two distinct solutions to the 4-queens puzzle as shown above

N given  
count num of ways to place N Queens in  
N x N matrix?





f(col) {

for (i=0 to n-1)  
 { if (fill -> ✓)  
 matrix[row][col] = 0  
 f(col+1)  
 }

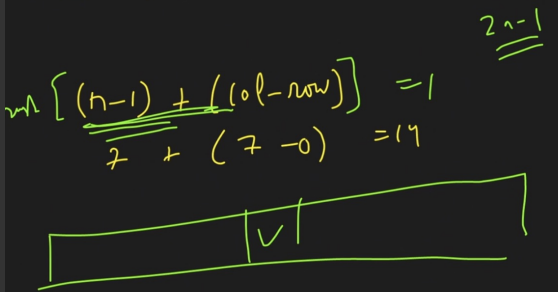
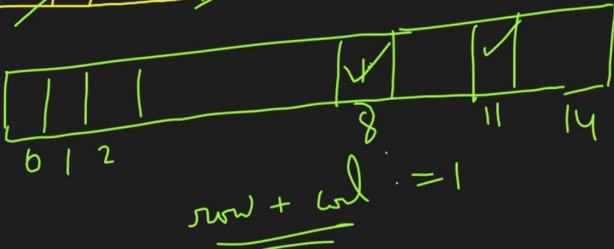
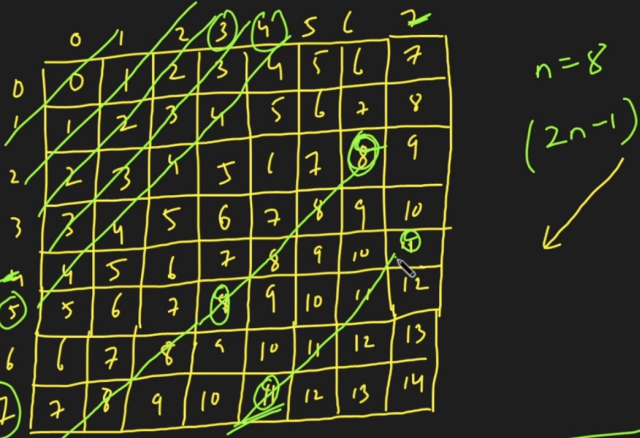
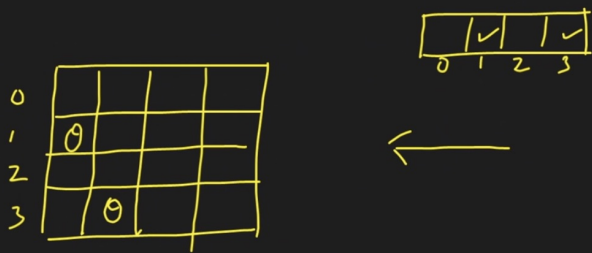
?

M2 -> optimal -> hashing

```

1 function isSafe1(row, col, board, n) {
2   let dupRow = row;
3   let dupCol = col;
4
5   while (row >= 0 && col >= 0) {
6     if (board[row][col] === "Q") return false;
7     row--;
8     col--;
9   }
10
11  col = dupCol;
12  row = dupRow;
13  while (col >= 0) {
14    if (board[row][col] === "Q") return false;
15    col--;
16  }
17
18  row = dupRow;
19  col = dupCol;
20  while (row < n && col >= 0) {
21    if (board[row][col] === "Q") return false;
22    row++;
23    col--;
24  }
25
26  return true;
27 }
28
29 function solve(col, board, ans, n) {
30   if (col === n) {
31     ans.push(...board);
32     return;
33   }
34
35   for (let row = 0; row < n; row++) {
36     if (isSafe1(row, col, board, n)) {
37       board[row] =
38         board[row].substring(0, col) +
39         "Q" +
40         board[row].substring(col + 1);
41       solve(col + 1, board, ans, n);
42       board[row] =
43         board[row].substring(0, col) +
44         "." +
45         board[row].substring(col + 1);
46     }
47   }
48 }
49
50 function solveNQueens(n) {
51   let ans = [];
52   let board = new Array(n).fill(".").repeat(n);
53   solve(0, board, ans, n);
54   return ans;
55 }

```



```

1 function solve(col, board, ans, leftRow, upperDiagonal, lowerDiagonal, n) {
2   if (col === n) {
3     ans.push([...board]);
4     return;
5   }
6
7   for (let row = 0; row < n; row++) {
8     if (
9       leftRow[row] === 0 &&
10      lowerDiagonal[row + col] === 0 &&
11      upperDiagonal[n - 1 + col - row] === 0
12    ) {
13      board[row] =
14        board[row].substring(0, col) +
15        "Q" +
16        board[row].substring(col + 1);
17      leftRow[row] = 1;
18      lowerDiagonal[row + col] = 1;
19      upperDiagonal[n - 1 + col - row] = 1;
20      solve(
21        col + 1,
22        board,
23        ans,
24        leftRow,
25        upperDiagonal,
26        lowerDiagonal,
27        n
28      );
29      board[row] =
30        board[row].substring(0, col) +
31        "." +
32        board[row].substring(col + 1);
33      leftRow[row] = 0;
34      lowerDiagonal[row + col] = 0;
35      upperDiagonal[n - 1 + col - row] = 0;
36    }
37  }
38 }
39
40 function solveNQueens(n) {
41   let ans = [];
42   let board = new Array(n).fill(".").repeat(n);
43
44   for (let i = 0; i < n; i++) {
45     board[i] = board[i];
46   }
47
48   let leftRow = new Array(n).fill(0);
49   let upperDiagonal = new Array(2 * n - 1).fill(0);
50   let lowerDiagonal = new Array(2 * n - 1).fill(0);
51
52   solve(0, board, ans, leftRow, upperDiagonal, lowerDiagonal, n);
53   return ans;
54 }

```

# lec17 sudoku solver

9x9 board → 3x3 boards

## Rules.

- 1. the digit 1-9 → once in any row
- 2. " " 1-9 → " " " col
- 3. " " 1-9 → " " " cell

## 37. Sudoku Solver

Hard 7.8K 207

Bloomberg Amazon Google

Write a program to solve a Sudoku puzzle by filling the empty cells.

A sudoku solution must satisfy all of the following rules:

- Each of the digits 1-9 must occur exactly once in each row.
  - Each of the digits 1-9 must occur exactly once in each column.
  - Each of the digits 1-9 must occur exactly once in each of the 9 3x3 sub-boxes of the grid.
- The '.' character indicates empty cells.

### Example 1:

5	3	.	.	7	.	.	.	.
6	.	.	1	9	5	.	.	.
.	9	8	.	.	.	.	6	.
8	.	.	.	6	.	.	.	3
4	.	.	8	.	3	.	.	1
7	.	.	.	2	.	.	.	6
.	6	.	.	.	2	8	.	.
.	.	.	4	1	9	.	.	5
.	.	.	.	8	.	.	7	9

```
Input: board = [
["5","3",".",".","7",".",".",".","."],
["6",".",".","1","9","5",".",".","."],
[".","9","8",".",".",".","6",".","."],
["8",".",".",".","6",".",".","3","."],
["4",".",".","8",".","3",".","6","."],
["7",".",".","2",".",".","6",".","."],
[".","6",".",".","2","8",".",".","."],
[".","4","1","9",".","5",".","7","9"]
]
```

```
Output: [
["5","3","4","6","7","8","9","1","2"],
["6","7","2","1","9","5","3","4","8"],
["1","9","8","3","4","2","5","6","7"],
["8","5","9","7","6","1","4","2","3"],
["4","2","6","8","5","3","7","9","1"],
["7","1","3","9","2","4","8","5","6"],
["9","6","1","5","3","7","2","8","4"],
["2","8","7","4","1","9","6","3","5"],
["3","4","5","2","8","6","1","7","9"]
]
```

Explanation: The input board is shown above and the only valid solution is shown below:

5	3	4	6	7	8	9	1	2
6	7	2	1	9	5	3	4	8
1	9	8	3	4	2	5	6	7
8	5	9	7	6	1	4	2	3
4	2	6	8	5	3	7	9	1
7	1	3	9	2	4	8	5	6
9	6	1	5	3	7	2	8	4
2	8	7	4	1	9	6	3	5
3	4	5	2	8	6	1	7	9



5	3		6	7	8	9		2
6	7	2	1	9	5	3	4	8
	9	8	3	4	2	5	6	7
8	5	9	7	6	1	4	2	3
4	2	6	8	5	3	7	9	1
7	1	3	9		4	8	5	6
9	6		5	3	7	2	8	4
2	8	7	4	1	9	6	3	5
3	4	5	2	8	6	1	7	9

5	3		6	7	8	9		2
6	7	2	1	9	5	3	4	8
	9	8	3	4	2	5	6	7
8	5	9	7	6	1	4	2	3
4	2	6	8	5	3	7	9	1
7	1	3	9		4	8	5	6
9	6		5	3	7	2	8	4
2	8	7	4	1	9	6	3	5
3	4	5	2	8	6	1	7	9

5	3		6	7	8	9		2
6	7	2	1	9	5	3	4	8
	9	8	3	4	2	5	6	7
8	5	9	7	6	1	4	2	3
4	2	6	8	5	3	7	9	1
7	1	3	9		4	8	5	6
9	6		5	3	7	2	8	4
2	8	7	4	1	9	6	3	5
3	4	5	2	8	6	1	7	9

try to fill up here  
 → not possible  
 so return false

after filling out  
 try out next  
 empty  
 go right then bottom

check to fill empty by trying  
 1 to 9

return 1 valid  
 sudoku

5	3		6	7	8	9		2
6	7	2	1	9	5	3	4	8
	9	8	3	4	2	5	6	7
8	5	9	7	6	1	4	2	3
4	2	6	8	5	3	7	9	1
7	1	3	9		4	8	5	6
9	6		5	3	7	2	8	4
2	8	7	4	1	9	6	3	5
3	4	5	2	8	6	1	7	9

5	3	4	6	7	8	9		2
6	7	2	1	9	5	3	4	8
	9	8	3	4	2	5	6	7
8	5	9	7	6	1	4	2	3
4	2	6	8	5	3	7	9	1
7	1	3	9		4	8	5	6
9	6		5	3	7	2	8	4
2	8	7	4	1	9	6	3	5
3	4	5	2	8	6	1	7	9

5	3	9	6	7	8	9	1	2
6	7	2	1	9	5	3	4	8
	9	8	3	4	2	5	6	7
8	5	9	7	6	1	4	2	3
4	2	6	8	5	3	7	9	1
7	1	3	9	4		8	5	6
9	6		5	3	7	2	8	4
2	8	7	4	1	9	6	3	5
3	4	5	2	8	6	1	7	9

M2 → optimal

check for row, col & cell in a single iteration

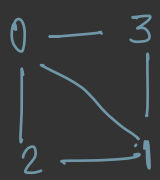
```

1 function solveSudoku(board) {
2   solve(board);
3 }
4
5 function solve(board) {
6   for (let i = 0; i < board.length; i++) {
7     for (let j = 0; j < board[0].length; j++) {
8       if (board[i][j] === ".") {
9         for (let c = 1; c <= 9; c++) {
10          if (isValid(board, i, j, c)) {
11            board[i][j] = c.toString();
12            if (solve(board)) {
13              return true;
14            } else {
15              board[i][j] = ".";
16            }
17          }
18        }
19        return false;
20      }
21    }
22  }
23  return true;
24 }
25
26 function isValid(board, row, col, c) {
27   for (let i = 0; i < 9; i++) {
28     if (board[i][col] === c.toString()) {
29       return false;
30     }
31     if (board[row][i] === c.toString()) {
32       return false;
33     }
34     if (
35       board[3 * Math.floor(row / 3) + Math.floor(i / 3)][
36         3 * Math.floor(col / 3) + (i % 3)
37       ] === c.toString()
38     ) {
39       return false;
40     }
41   }
42   return true;
43 }

```

lec 18. M-coloring problem

at most M colors M=3



{(0,3), (3,1), (1,2), (2,0), (0,1)}

**M-Coloring Problem**

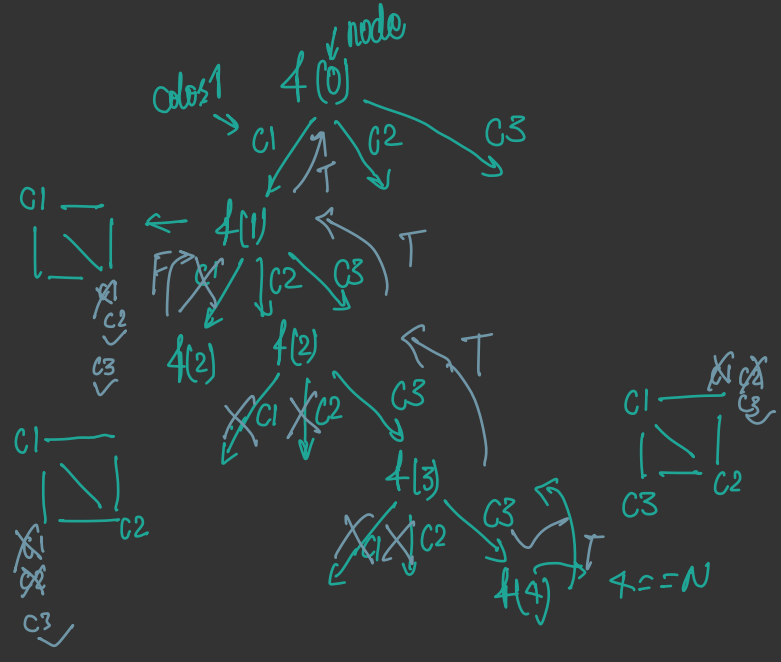
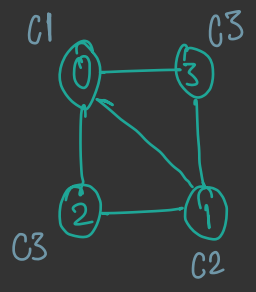
Medium Accuracy: 33.66% Submissions: 6921 Points: 4

Given an undirected graph and an integer **M**. The task is to determine if the graph can be colored with at most **M** colors such that no two adjacent vertices of the graph are colored with the same color. Here coloring of a graph means the assignment of colors to all vertices. Print 1 if it is possible to colour vertices and 0 otherwise.

**Example 1:**

**Input:**  
 N = 4  
 M = 3  
 E = 5  
 Edges [] = {(1,2), (2,3), (3,4), (4,1), (1,3)}

**Output:** 1  
**Explanation:** It is possible to colour the given graph using 3 colours.



f(node) {

if (node == N) return T;

for (col = 1 to m)

{ if (possible ✓)

{ color[node] = col;

if (f(node+1) == T)

return T;

color[node] = 0;

}

return F;

}

TC  $\rightarrow O(N^M) \rightarrow$  try m colors at each node

SC  $\rightarrow O(N) + O(N)$

recursion

```

1 function isSafe(node, G, color, n, col) {
2   for (let item of G[node]) {
3     if (color[item] === col) return false;
4   }
5   return true;
6 }
7
8 function solveGraph(node, G, color, n, m) {
9   if (node === n) return true;
10  for (let i = 1; i <= m; i++) {
11    if (isSafe(node, G, color, n, i)) {
12      color[node] = i;
13      if (solveGraph(node + 1, G, color, n, m)) return true;
14      color[node] = 0;
15    }
16  }
17  return false;
18 }
19
20 function graphColoring(G, color, i, m) {
21  const n = G.length;
22  if (solveGraph(i, G, color, n, m)) return true;
23  return false;
24 }

```

### lec 19. Palindrome Partitioning.

aabb  $\rightarrow$  {a, a, b, b}

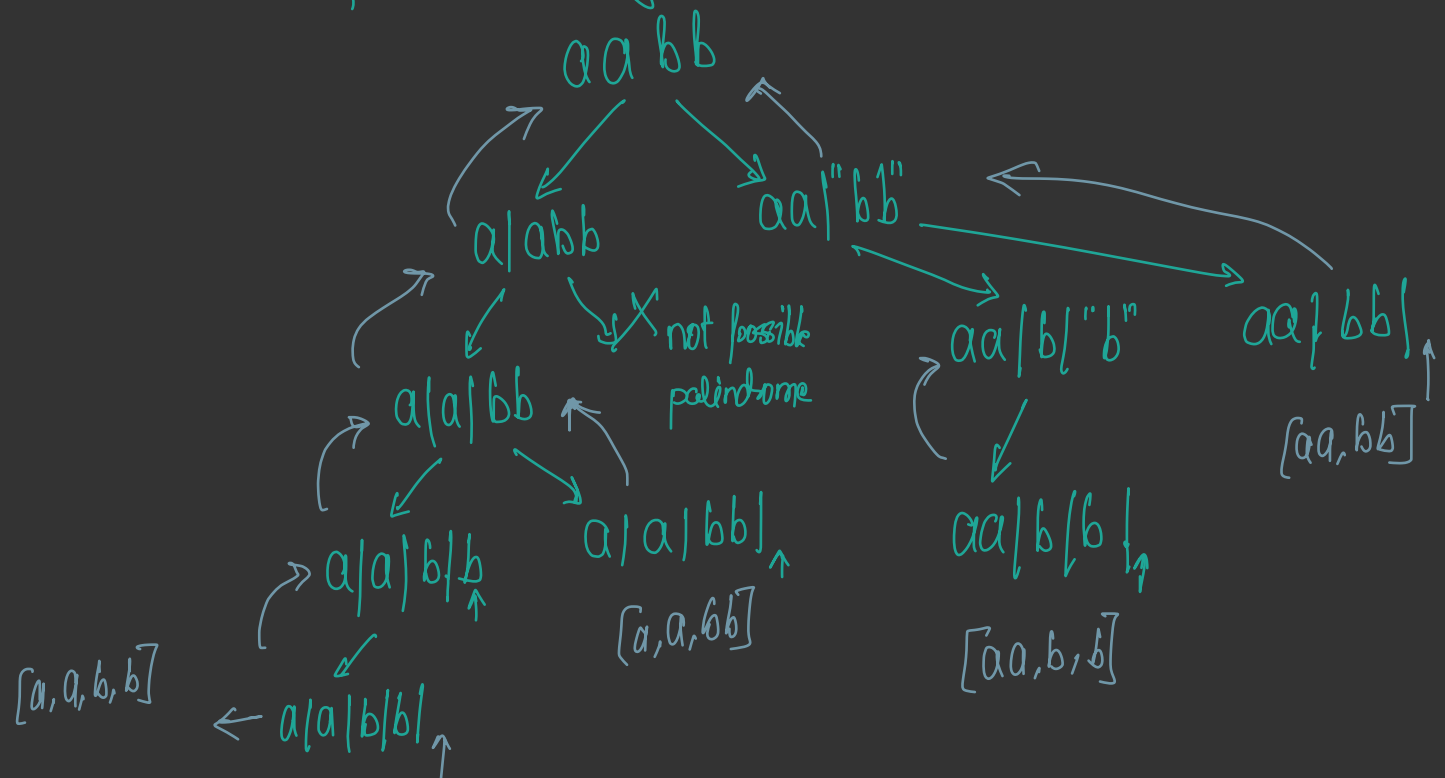
{a, a, bb}

{aa, bb}

{aa, bb}

### M1 brute force.

we can do partition only where left part is a palindrome



### 131. Palindrome Partitioning

Medium

10.6K

335

☆

🔄

Bloomberg

Amazon

Google

...

Given a string `s`, partition `s` such that every substring of the partition is a **palindrome**. Return all possible palindrome partitioning of `s`.

#### Example 1:

Input: `s = "aab"`

Output: `[["a","a","b"],["aa","b"]]`

```

4(ind) {
  if(ind == n) {
    ans.push(path)
    return;
  }
  for(i = ind; i < n; i++) {
    if(isPalindrome(s, ind, i)) {
      path.push(s.substring(ind, i+1))
      4(i+1)
      path.pop()
    }
  }
}

```

```

1 const isPalindrome = (s, start, end) => {
2   while (start <= end) {
3     if (s.charAt(start++) !== s.charAt(end--)) return false;
4   }
5   return true;
6 };
7
8 const func = (index, s, path, res) => {
9   if (index === s.length) {
10    res.push([...path]);
11    return;
12  }
13
14  for (let i = index; i < s.length; ++i) {
15    if (isPalindrome(s, index, i)) {
16      path.push(s.substring(index, i + 1));
17      func(i + 1, s, path, res);
18      path.pop();
19    }
20  }
21 };
22
23 const partition = (s) => {
24   const res = [];
25   const path = [];
26   func(0, s, path, res);
27   return res;
28 };

```

lec20 => Rat in a maze

this color => backtrack

**Rat in a Maze Problem - I** Medium Accuracy: 37.73% Submissions: 51335 Points: 4

Consider a rat placed at (0, 0) in a square matrix of order N \* N. It has to reach the destination at (N - 1, N - 1). Find all possible paths that the rat can take to reach from source to destination. The directions in which the rat can move are 'U'(up), 'D'(down), 'L'(left), 'R'(right). Value 0 at a cell in the matrix represents that it is blocked and rat cannot move to it while value 1 at a cell in the matrix represents that rat can be travel through it.

**Note:** In a path, no cell can be visited more than one time.

**Example 1:**

**Input:**  
 N = 4  
 m[][] = {{1, 0, 0, 0},  
 {1, 1, 0, 1},  
 {1, 1, 0, 0},  
 {0, 1, 1, 1}}

**Output:**  
 DDRDRR DRDRR

	0	1	2	3
0	1	0	0	0
1	0	1	0	1
2	0	0	0	0
3	0	0	0	1

✓			
✓			
✗	✗		
	✗	✗	

vis

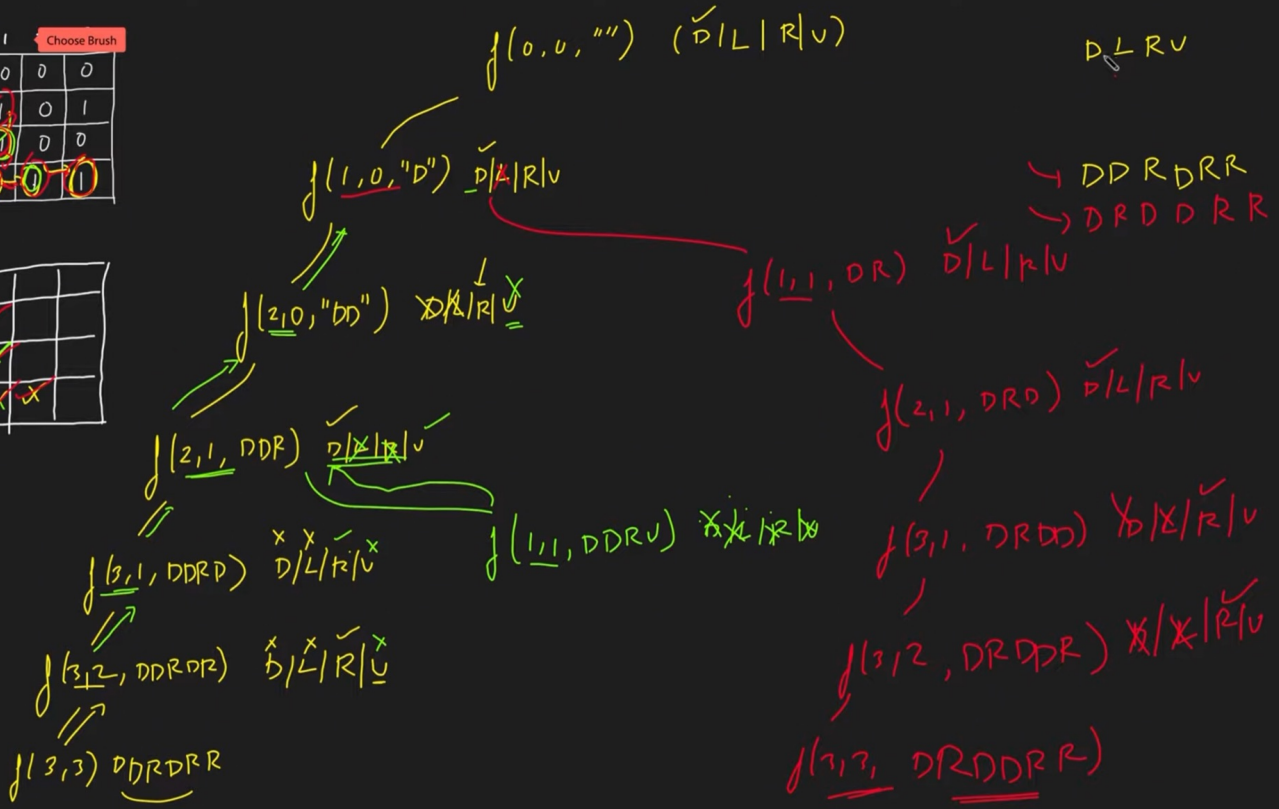
$f(0,0, "")$  (D|L|R|U) DLRU  
 $f(1,0, "D")$  D|L|R|U DDRDRR  
 $f(2,0, "DD")$  D|L|R|U  
 $f(2,1, "DDR")$  D|L|R|U  
 $f(3,1, "DDRDRR)$  D|L|R|U  
 $f(3,2, "DDRDRR)$  D|L|R|U  
 $f(3,3, "DDRDRR)$



	0	1	Choose Brush
0	0	0	0
1	0	1	0
2	0	0	0
3	0	0	0

✓			
✓	✓		
✓	✓	✓	
		✓	✓

vis



TC  $\rightarrow 4^{N \times M}$

SC  $\rightarrow O(M \times N)$

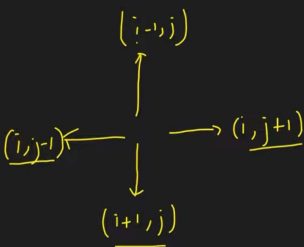
↑ RECURSION

{  
 down  
 f(i)  
 left  
 f(i)  
 right  
 f(i)  
 up  
 f(i)  
 }

```

1  const solve = (i, j, a, n, ans, move, vis) => {
2      if (i === n - 1 && j === n - 1) {
3          ans.push(move);
4          return;
5      }
6
7      // downward
8      if (i + 1 < n && !vis[i + 1][j] && a[i + 1][j] === 1) {
9          vis[i][j] = 1;
10         solve(i + 1, j, a, n, ans, move + "D", vis);
11         vis[i][j] = 0;
12     }
13
14     // left
15     if (j - 1 >= 0 && !vis[i][j - 1] && a[i][j - 1] === 1) {
16         vis[i][j] = 1;
17         solve(i, j - 1, a, n, ans, move + "L", vis);
18         vis[i][j] = 0;
19     }
20
21     // right
22     if (j + 1 < n && !vis[i][j + 1] && a[i][j + 1] === 1) {
23         vis[i][j] = 1;
24         solve(i, j + 1, a, n, ans, move + "R", vis);
25         vis[i][j] = 0;
26     }
27
28     // upward
29     if (i - 1 >= 0 && !vis[i - 1][j] && a[i - 1][j] === 1) {
30         vis[i][j] = 1;
31         solve(i - 1, j, a, n, ans, move + "U", vis);
32         vis[i][j] = 0;
33     }
34 };
35
36 const findPath = (m, n) => {
37     const ans = [];
38     const vis = Array.from({ length: n }, () => Array(n).fill(0));
39     if (m[0][0] === 1) {
40         solve(0, 0, m, n, ans, "", vis);
41     }
42     return ans;
43 };

```



D | L | R | U

$di[] = +1 \quad +0 \quad +0 \quad -1$   
 $dj[] = +0 \quad -1 \quad +1 \quad +0$

$i + di[md]$   
 $j + dj[md]$

```

1 const solve = (i, j, a, n, ans, move, vis, di, dj) => {
2   if (i === n - 1 && j === n - 1) {
3     ans.push(move);
4     return;
5   }
6
7   const dir = "DLRU";
8   for (let ind = 0; ind < 4; ind++) {
9     const nexti = i + di[ind];
10    const nextj = j + dj[ind];
11
12    if (
13      nexti >= 0 &&
14      nextj >= 0 &&
15      nexti < n &&
16      nextj < n &&
17      !vis[nexti][nextj] &&
18      a[nexti][nextj] === 1
19    ) {
20      vis[i][j] = 1;
21      solve(nexti, nextj, a, n, ans, move + dir[ind], vis, di, dj);
22      vis[i][j] = 0;
23    }
24  }
25 };
26
27 const findPath = (m, n) => {
28   const ans = [];
29   const vis = Array.from({ length: n }, () => Array(n).fill(0));
30   const di = [1, 0, 0, -1];
31   const dj = [0, -1, 1, 0];
32
33   if (m[0][0] === 1) {
34     solve(0, 0, m, n, ans, "", vis, di, dj);
35   }
36
37   return ans;
38 };

```

loc 21 kth permutation sequence

M1 bruteforce

→ generate all permutations and store it in a data structure ans

→ sort ans

→ and return (k-1) index item

TC →  $N! \times N + O(N! \times N \log N)$   
 ↑ deep copy of ds

M2 optimal soln

n=4 k=17

total permutations =  $4! = 4 \times 3 \times 2 = 24$

[1, 2, 3, 4]  
 0 1 2 3

1 2 3 4 → 0th permutation

4 3 2 1 → 23th

**60. Permutation Sequence**

Hard 2081 360 Add to List Share

The set  $\{1, 2, 3, \dots, n\}$  contains a total of  $n!$  unique permutations.

By listing and labeling all of the permutations in order, we get the following sequence for  $n = 3$ :

- "123"
- "132"
- "213"
- "231"
- "312"
- "321"

Given  $n$  and  $k$ , return the  $k^{\text{th}}$  permutation sequence.

**Example 1:**  
 Input:  $n = 3, k = 3$   
 Output: "213"

**Example 2:**  
 Input:  $n = 4, k = 9$   
 Output: "2314"

**Example 3:**  
 Input:  $n = 3, k = 1$   
 Output: "123"

if starts with

$1 + (2, 3, 4) ] 6 \quad (0-5)$   
 $2 + (1, 3, 4) ] 6 \quad (6-11)$   
 $3 + (1, 2, 4) ] 6 \quad (12-17)$   
 $4 + (1, 2, 3) ] 6 \quad (18-23)$

24

3

return  $k-1 = 16^{\text{th}} \Rightarrow 16/6 = 2$   
 in each set  $\uparrow$  at index 2  $\Rightarrow 3$

1 2 3 4  
0 1 2 3

$16/6 = 4 \Rightarrow 4^{\text{th}}$  sequence in (12-17)

$[1, 2, 4]$ ,  $k=4$   
 $4/2 \Rightarrow 2$   
 $4/2 \Rightarrow 0$

1 + {2,4} } 2 (0-1)  
 2 + {1,4} } 2 (2-3)  
 4 + {1,2} } 2 (4-5)  
 6

3 4 1 2  
 remaining  $\uparrow$

0th sequence generated by {1,2}  
 $\{1, 2\}$   $0/1 = 0$   
 1 + {2} } 1 (0-0)  
 2 + {1} } 1 (1-1)  
 2

$n=4, k=17 = 16^{\text{th}}$   $k=16/6$

$(4!) = 24$   
 $3! = 6$   
 $2! = 2$   
 $1! = 1$

$\{1, 2, 3, 4\}$ ,  $k = (17)/2 = 8$   
 $3! = 6$   $4/2 = 0$

$\{1, 2\}$   $k = 0/1 = 0$   
 $2! = 2$   $k = 0/1 = 0$

$\{2\}$   $k=0$  (3 4 1 2) 16th

1 6  
 6  
 0  $\leftarrow 16/6 = 2$   
 6  
 2 0-1  
 2 2-3  
 2 4-5  
 1

TC  $\Rightarrow O(N) \times O(N)$   
 $\Rightarrow O(N^2)$   
 SC  $\rightarrow O(N)$

```

1  const getPermutation = (n, k) => {
2  let fact = 1;
3  const numbers = [];
4
5  for (let i = 1; i < n; i++) {
6  fact = fact * i;
7  numbers.push(i);
8  }
9
10 numbers.push(n);
11 let ans = "";
12
13 k = k - 1;
14
15 while (true) {
16 ans = ans + numbers[Math.floor(k / fact)].toString();
17 numbers.splice(Math.floor(k / fact), 1);
18
19 if (numbers.length === 0) {
20 break;
21 }
22
23 k = k % fact;
24 fact = Math.floor(fact / numbers.length);
25 }
26
27 return ans;
28 };
    
```